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## LARSON—OPER 731—CLASSROOM WORKSHEET 06 Linear Programs

## Concepts

- (Sec. 2.1) infeasible linear program, certificate of infeasibility, certificate of unboundedness, certificate of optimality.
- (Sec. 2.2) SEF, *slack variables*.
- (Sec. 2.4) basis, basic variable, nonbasic variable, basic solution, basic feasible solution, canonical form.

## **Bases and Canonical Forms**

- 1. If A is  $m \times n$  and  $j \in \{1, \ldots, n\}$ , what is  $A_j$ ? If  $J \subseteq \{1, \ldots, n\}$ , what is  $A_J$ ?
- 2. If  $B \subseteq \{1, \ldots, n\}$  when is B a basis for A? What can we say about  $A_B$ ? What is N?
- 3. What are *basic* variables  $x_j$ ? What are *nonbasic* variables?
- 4. What is a *basic solution* of Ax = b for a basis B? (What is a *basic feasible solution*?)
- 5. What is the *canonical form* for an LP with respect to a basis B?
- 6. Consider the following LP in SEF:

max (1, -2, 0, 1, 3)xsubject to

$$\begin{pmatrix} 1 & -1 & 2 & -1 & 0 \\ 2 & 0 & 1 & -1 & 1 \end{pmatrix} x = \begin{pmatrix} 1 \\ -1 \end{pmatrix}$$
  
  $x \ge 0.$ 

- 7. Check that  $\{1, 4\}$  is a basis.
- 8. Find an equivalent LP in canonical form.

## Simplex Method Steps

$$max\{c^Tx + \bar{z} : Ax = b, x \ge \mathbb{O}\}\$$

- (a) Check that Ax = b has linearly independent rows (using Gaussian elimination).
- (b) Check that Ax = b is feasible by solving the system (using Gaussian elimination) or discovering that the system is infeasible. Assume A is  $m \times n$ . You'll get m leading 1's. Let B be the indices of those 1's columns. The initial x can be taken to be the solution vector with 0's for all non-B components (so  $x_N = \mathbb{O}$ ).
- (c) We have initial basis B and initial basic feasible solution  $\bar{x}$ . Now write Ax = b in canonical form with respect to B:
  - i. On the first pass  $A_B = I$  (due to Gaussian elimination). Otherwise leftmultiply Ax = b by  $A_B^{-1}$ .
  - ii. Let  $y = (y_1, ..., y_m)^T$ . Then  $y^T A x = y^T b$ , and  $y^T b y^T A x = \mathbb{O}$ . Add this to  $c^T x + \overline{z}$  to get and updated objective function.
  - iii. We want y so that our updated objective function has  $c_B = 0$ . The new objective function is  $(c^T x + \bar{z}) + (y^T b y^T A x)$ . Solve to get:  $c_B^T x_B - y^T A_B x_B = \mathbb{O}$ .
- (d) Let c' be the x coefficients in the updated objective function (with  $c'_B = \mathbb{O}$ ). Are any c' entries positive? If not we're done ( $\hat{x}$  is optimal). Let k be the smallest index with  $c'_k > 0$  (k is necessarily in N). We want to increase  $\bar{x}_k$  (which is necessarily 0) to t.
- (e) We'll keep all other  $\bar{x}$  components for indices in N equal to 0. So,  $b = A_B x_B + x_k A_k = x_B + t A_k$ . So  $x_B = b - t A_k \ge \mathbb{O}$ , and  $b \ge t A_k$ .
- (f) For each component  $b_i$  of b, we must have  $\frac{b_i}{A_{ik}}$  (when  $A_{ik} \neq 0$ ). Let  $t = \min\{\frac{b_i}{A_{ik}} : i \in B\}.$
- (g) Let r be the smallest index where  $t = \frac{b_i}{A_{ik}}$ .
- (h) Let the updated feasible solution be  $\bar{x}'$  with  $\bar{x}' = b tA_k$ ,  $\bar{x}'_k = t$  (and  $\bar{x}'_i = 0$  for  $i \in N \setminus \{k\}$ ).
- (i) Remove r from B and add k. So the new basis is  $B' = B \setminus \{r\} \cup \{k\}$ .
- (j) Add r to N and remove k. So the new nonbasis is  $N' = N \setminus \{k\} \cup \{r\}$ .
- (k) We now have new basis B', nonbasis N', basic feasible  $\bar{x}'$  (with  $\bar{x}'_N = \mathbb{O}$ ), and objective function with x coefficients c' (with  $c'_B = \mathbb{O}$ ).
- 9. Solve.
  - **1** Consider the LP problem max { $c^{T}x$  :  $Ax = b, x \ge 0$ }, where

$$A = \begin{pmatrix} 1 & 2 & -2 & 0 \\ 0 & 1 & 3 & 1 \end{pmatrix} \qquad b = \begin{pmatrix} 2 \\ 5 \end{pmatrix} \qquad c = \begin{pmatrix} 0 \\ 3 \\ 1 \\ 0 \end{pmatrix}.$$

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(a) Beginning with the basis  $B = \{1, 4\}$ , solve the problem with the simplex method. At each step, choose the entering variable and leaving variable by Bland's rule.